

Middleware and Virtualization Management

# Resource Co-Allocation for Large-Scale Distributed Environments

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### Motivation (1)

- Co-allocation of resources: allocation of multiple resources within the same time window
- Emergence of new paradigms
  - -On Demand computing (Amazon EC2, SalesForce, IBMCloud)
- Requirements
  - -QoS/SLA support
  - -Efficiency
  - -Scalability
- Emergence of new applications that capitalized on the availability of distributed computing to perform tasks with spatial and temporal dependencies (MapReduce/financial apps.)



# Motivation (2)

#### Applications

- -Virtual Computing Lab (VCL)
- –MapReduce framework (Hadoop)
- -Grid lambda scheduling
- -Workflow scheduling

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# Background

- Naïve Approach: A co-allocation request can be treated as a group of sequential scheduling requests
  - Inappropriate for time sensitive applications

#### Batch scheduling

- Resource driven: optimizing for system performance
  - Limited support for QoS by means of backfilling and priorities
- Job driven: optimizing for application performance

#### Advance reservations

- -QoS provisioning
- -Workflow support
- -Multiple drawbacks



### Goals

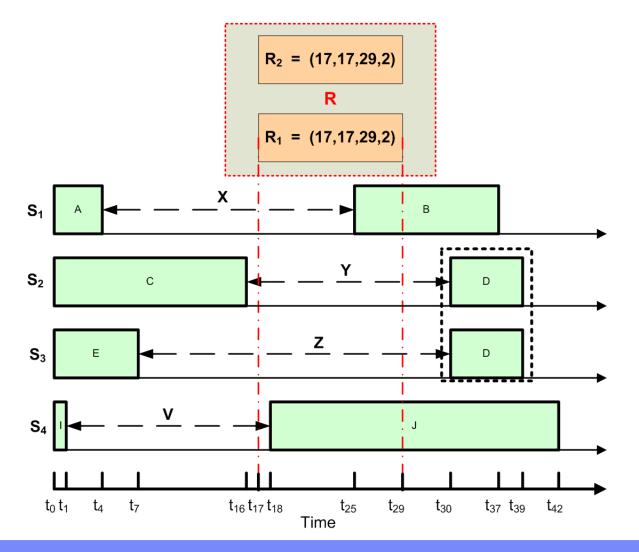
- Providing users with time guarantees by scheduling jobs as they arrive without promoting resource fragmentation
- Allowing better scheduling decisions by keeping look ahead until the horizon of the schedule in a way that is *efficient*

#### Contributions

- A co-allocation scheduling algorithm
  - Effective in co-allocating resources and provides support for advance reservations and range search
- Range search
  - Ability of the system to find a set of resources available within a given time window
  - Enable selection and scheduling algorithms that are application specific
- Efficient data structure to organize resource availability
  - Leading to the design of an algorithm that allows a single search operation to identify all required resources *efficiently*



### **Problem Description**



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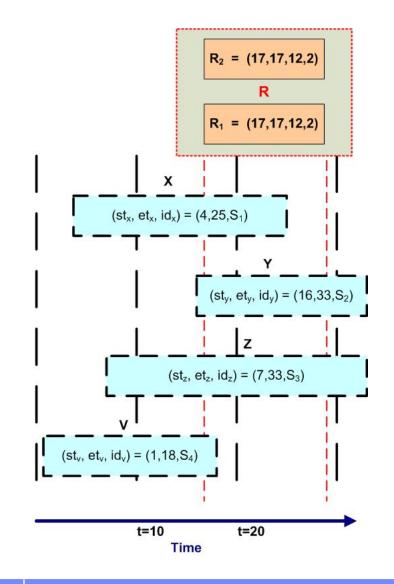
# System Model

#### We consider the following settings:

- -Scheduler S
- -N servers
- -Reservation request *r* requires service
- -Request (**q**<sub>r</sub>,**s**<sub>r</sub>,**I**<sub>r</sub>,**n**<sub>r</sub>)
  - **q**<sub>r</sub> request time
  - s<sub>r</sub> earliest time the reservation is needed
  - I<sub>r</sub> temporal size of the request (duration)
  - n<sub>r</sub> spatial size of the request (no. or servers)
- -Idle period (st, et, id,)
  - st<sub>i</sub> starting time
  - et<sub>i</sub> ending time
  - id, server offering the idle period



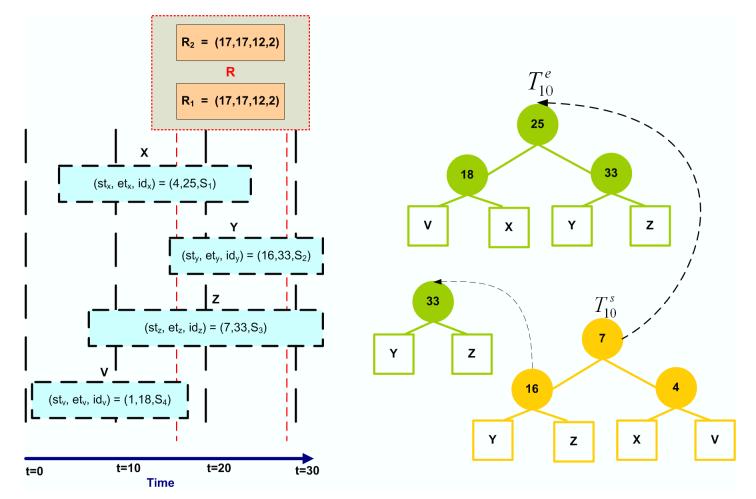
### **Data Structure and Algorithm**



- Time space is partitioned into time slots of equal length
- Idle periods are stored in each time slot they span over
- Algorithms searches only into the time slot containing s<sub>r</sub>
- Upon failure to schedule: s<sub>r</sub> = s<sub>r</sub> + Δ<sub>t</sub>
- Honor atomicity of the request by means of temporal counters
- Number of idle periods per time slot can be bounded to N if time slot size is set to the minimum temporal size



### **Data Structure**



Two feasibility criterion:  $\underline{st_i < st_x}$  and  $\underline{et_i > et_x}$ 



# **Performance Evaluation**

#### Real workloads drive simulations [ParArch]

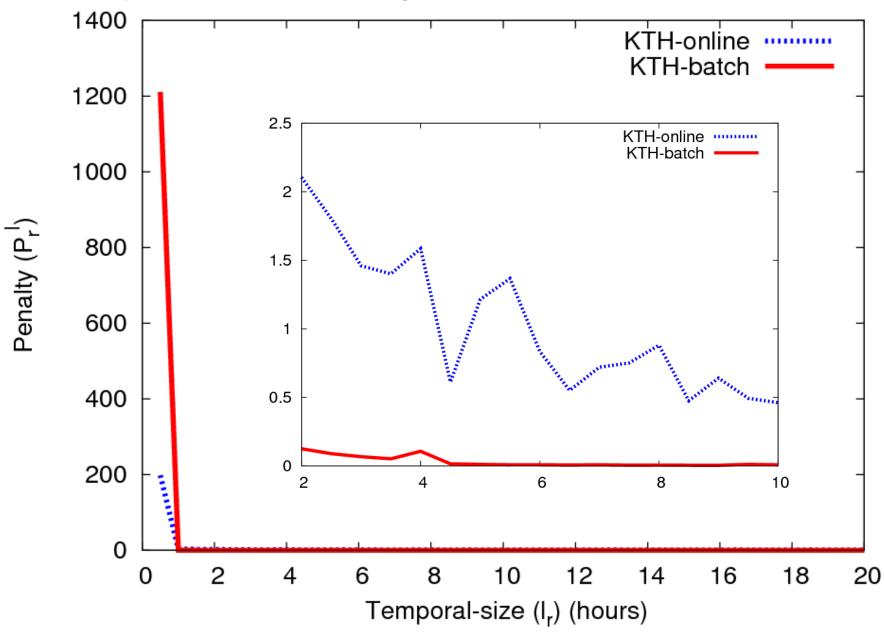
Workload	No. of processors	No. of jobs	Avge. length (hrs)	Avge. spatial size
СТС	512	39,734	5.82	9.48
КТН	128	28,481	2.46	7.67
HPC2N	240	202,825	4.72	6.56

Two experiments

- -Comparison to batch scheduling
- -Impact on performance of advance advance-reservations



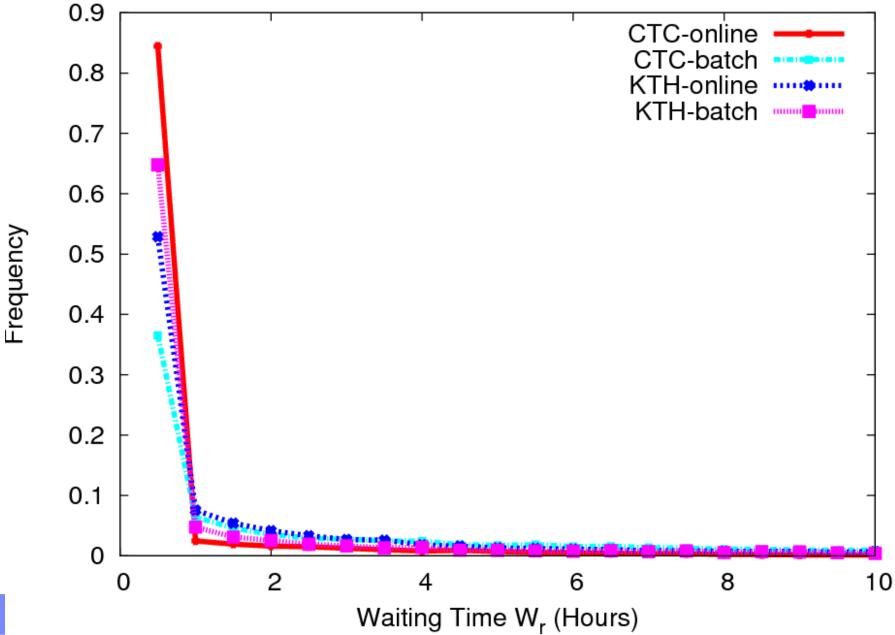
### **Temporal-size Penalty**



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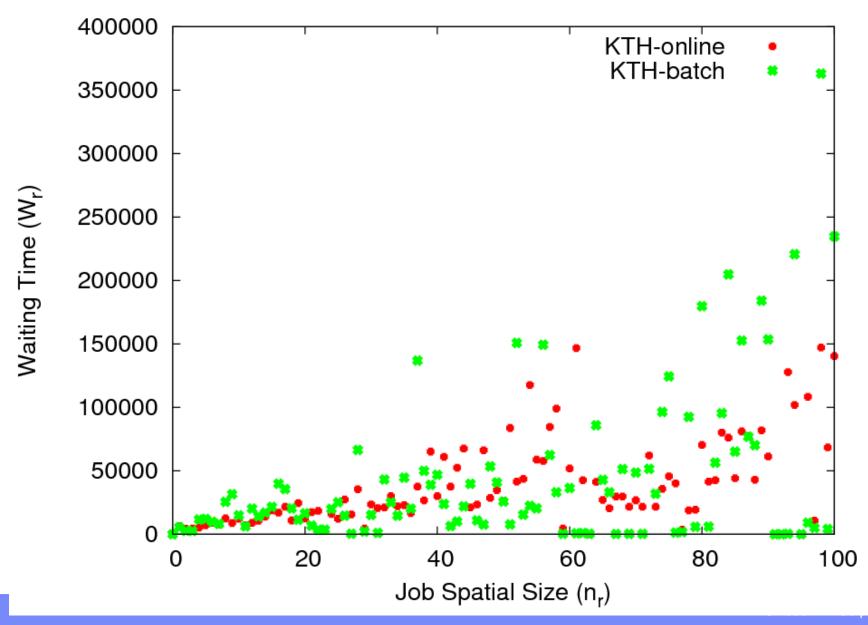


### Waiting Time Distribution

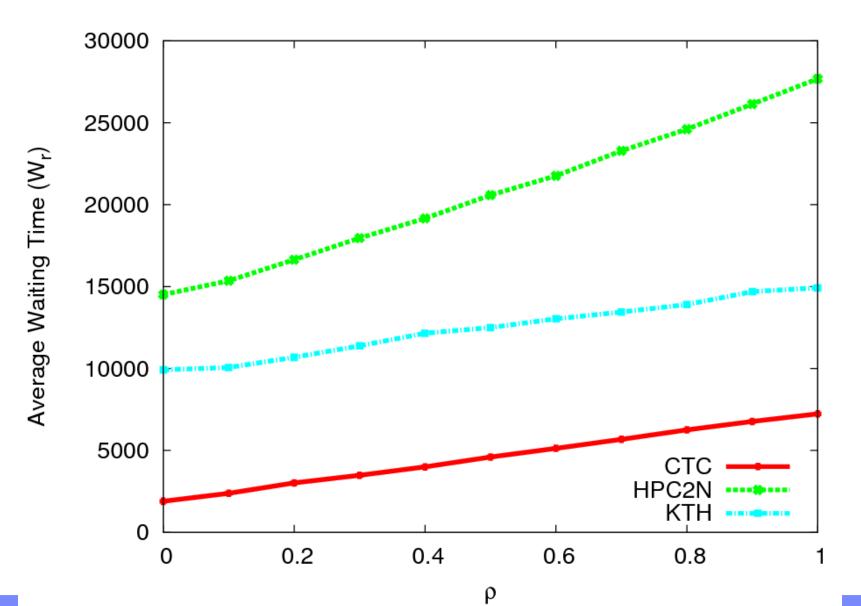


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#### Waiting time distribution as a function of spatial size

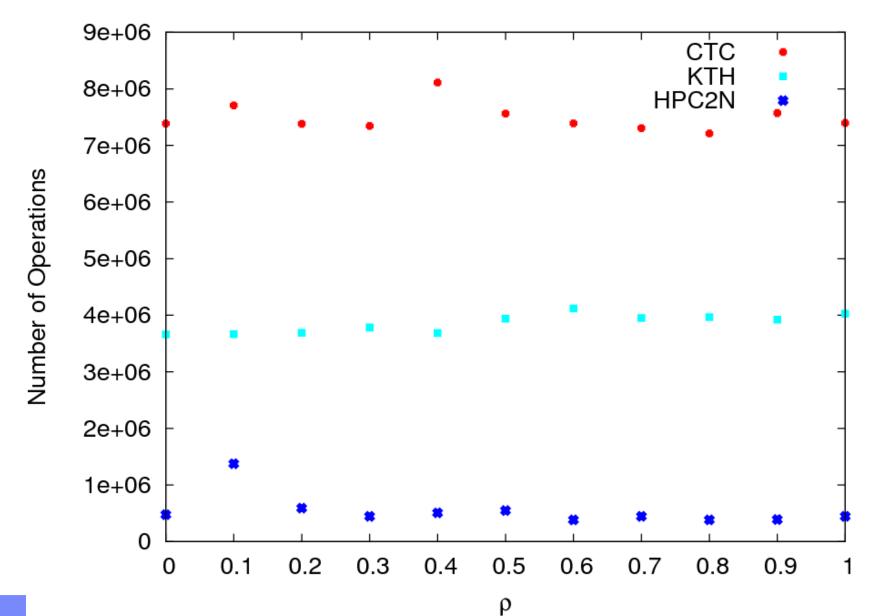


### Avg Waiting Time vs. fraction of advance reservations ( $\rho$ )





### Number of operations vs. fraction of advance reservations ( $\rho$ )





#### Number of retrials vs. spatial size

Workload/n	(0:50]	(50:100]	(100:150]	(150:200]	(350:400]
CTC (No. of retrials)	2.96	5.34	7.22	13.25	127.44
KTH (No. of retrials)	10.27	60	120		

- Larger spatial size distribution results in larger number of attempts.
- Temporal size distribution of KTH shows large proportion of small jobs.



# **Discussion of Results**

- Our algorithm can efficiently co-allocate resources while supporting advance reservations
- Online advance reservations mechanisms might offer a better solution to the problem of co-allocating resources as compared to conventional batch scheduling
- Our work can be easily extended to support deadlines



### **Future Work**

- Implement the co-allocation algorithm proposed in the context of
  - -Hadoop
  - -End-to-end path problem in Grid lambda scheduling
- Impact of workload characteristics on system/user performance
- Uncertainty of completion times



### Thank you!

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